## Computational Methods for Quantum Many-Body Physics

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## Logistical details

- Class timings: Wednesdays 15:00-17:00
- Venue: Emmy Noether seminar room and online
- Zoom details:
- Meeting ID: 87633971378
- Passcode: 171723
- Evaluation: Assignments (will be posted on the Moodle page)
- Contacts:
- Sthitadhi Roy [sthitadhi.roy@icts.res.in]
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- All course material (slides, lecture notes, video recordings, and assignments) can be found at https://courses.icts.res.in/course/view.php?id=82


## Purpose of this course

- obvious, introduce some the most commonly used numerical methods in many-body physics
- introduce how and why the methods work? understanding the algorithms behind methods
- how to extract useful and interesting physics from the numerical methods
- understand the physics behind the algorithms: If and how the algorithms work also encodes a lot of physics of the system under consideration
- At a practical level, lots of very well structured libraries available such as
- ALPS [https://alpscore.org]
- QuSpin [http://quspin.github.io/QuSpin/] [SciPost Phys. 2, 003 (2017), SciPost Phys. 7, 020 (2019)]
- TenPY [https://tenpy.readthedocs.io/en/latest/][SciPost Phys. Lect. Notes 5 (2018)]
- iTensor [https://itensor.org][sciPost Phys. Codebases 4 (2022)]
- many many others...

Help understand how do the underlying algorithms in these packages so that they are no longer black boxes

## Course Plan

- Lectures 1-3: Exact diagonalisation
- representing Hamiltonians as sparse matrices
- Lanczos algorithms for diagonalisation
- Shift-invert and Polynomially filtered exact diagonalisation
- time-evolution using ED
- Lectures 4-7: Classical and Quantum Monte Carlo:
- Basic principles of Monte Carlo algorithms: importance sampling, detailed balance, autocorrelation timescales, error analysis
- Illustrating classical Monte Carlo using the 2D Ising model: local, worm and cluster type algorithms
- Some other useful tricks: parallel tempering, overrelaxation etc
- Illustrating quantum Monte Carlo (QMC) using the 2D $S=1 / 2$ Heisenberg antiferromagnet and the 2D J-Q model: Stochastic series expansion (SSE) QMC and its implementation
- Introduction to sign problem (time permitting)
- Lecture 8: Time-evolution of quantum systems
- (truncated) Krylov space methods
- Kernel polynomial methods
- Lectures 9-10: Tensor Network Methods
- matrix product states
- matrix product operators
- introduction to DMRG and tDMRG algorithms (time permitting)

Models: typically disordered, interacting quantum many-body Hamiltonians defined on regular, hierarchical or random lattices

- Spin models

$$
H=\sum_{\langle i, j\rangle} \sum_{\mu, \nu=x, y, z} J_{i j}^{\mu \nu} \hat{S}_{i}^{\mu} \hat{S}_{j}^{\nu}+\sum_{i} \sum_{\mu=x, y, z} h_{i}^{\mu} \hat{S}_{i}^{\mu}
$$

- Fermi-Hubbard type models

$$
H=-t \sum_{\langle i, j\rangle, \sigma} c_{i \sigma}^{\dagger} c_{j \sigma}+\sum_{i, \sigma} \epsilon_{i, \sigma} c_{i \sigma}^{\dagger} c_{i \sigma}+U \sum_{i} n_{i \uparrow} n_{i \downarrow}
$$

- Bose-Hubbard type models

$$
H=-t \sum_{\langle i, j\rangle} b_{i}^{\dagger} b_{j}+\sum_{i} \epsilon_{i, \sigma} b_{i}^{\dagger} b_{i}+\sum_{i, j} V_{i j} n_{i} n_{j}
$$



- many more: Bose-Fermi mixtures, $t-J$ models, electrons coupled to phonons, open quantum systems ...
- Eigenvalue spectrum
- analysis of the gap across quantum phase transition
- spectral correlations: ergodic, chaotic or integrable, localised
- density of states, zero modes
- Correlation functions in ground states/eigenstates
- correlation lengths: diverges across transition, long-range order
- order parameters; scaling across QPTs
- entanglement structure in the states
- Thermodynamic properties
- partition functions and thermodynamic potentials
- response functions, specific heat, susceptibilities
- Non-equilibrium dynamics
- dynamical response, structure factors
- non-equilibrium transport
- quantum chaos or lack thereof


Schmitt et al., Sci. Adv. 2022; Sandvik's lecture note, Karrash et al. PRB 2014

## Exact Diagonalisation

1. Construct and enumerate the basis states

- label the sites of your lattice
- enumerate the basis states $\Rightarrow$ assign an unique integer label to each basis state
- pick out the subset of basis states allowed by symmetries/conservation laws


## 2. Construct the Hamiltonian as a matrix

- identify the diagonal elements of the Hamiltonian $\Rightarrow$ these just associate numbers to each basis state
- identify the off-diagonal terms in the Hamiltonian $\Rightarrow$ for each term identify the set of basis states every basis state is connected to under the action of the term

3. Diagonalise the Hamiltonian $\Rightarrow$ extract the required eigenvalues and eigenvectors
4. Compute observabes

- express the operators of interest as matrices in the constructed basis
- expectation values as matrix-vector multiplications


## Single particle hopping on some random lattice

$$
H=-t \sum_{\langle i j\rangle} c_{i}^{\dagger} c_{j}+\sum_{i} \epsilon_{i} c_{i}^{\dagger} c_{i}
$$

- real-space basis a natural choice

- $|2\rangle=[0,0,1,0,0,0]^{\top}$ is a basis state localised on site 2 of the lattice
- onsite potentials $\left\{\epsilon_{i}\right\}$ form the diagonal elements
- hopping between sites connected by lines constitute the off-diagonal elements

$$
H=\left(\begin{array}{cccccc}
\epsilon_{0} & 1 & 1 & 0 & 1 & 0 \\
1 & \epsilon_{1} & 1 & 0 & 0 & 1 \\
1 & 1 & \epsilon_{2} & 1 & 0 & 0 \\
0 & 0 & 1 & \epsilon_{3} & 1 & 1 \\
1 & 0 & 0 & 1 & \epsilon_{4} & 1 \\
0 & 1 & 0 & 1 & 1 & \epsilon_{5} .
\end{array}\right)
$$

- Interacting spin Hamiltonian on a 1D chain

$$
H=J \sum_{i=0}^{L-1}\left[\sigma_{i}^{x} \sigma_{i+1}^{x}+\sigma_{i}^{y} \sigma_{i+1}^{y}+\sigma_{i}^{z} \sigma_{i+1}^{z}\right]+\sum_{i} h_{i} \sigma_{i}^{z}
$$

- Computational basis: classical configurations of $\sigma^{2}$-product states
- Enumeration of the basis states: spin-configurations $\rightarrow$ binary strings $\rightarrow$ integers

| Configuration | String | Integer |
| :---: | :---: | :---: |
| $\uparrow \uparrow \uparrow \uparrow$ | 0000 | 0 |
| $\uparrow \uparrow \uparrow \downarrow$ | 0001 | 1 |
| $\uparrow \uparrow \downarrow \uparrow$ | 0010 | 2 |
| $\uparrow \uparrow \downarrow \downarrow$ | 0011 | 3 |
| $\uparrow \downarrow \uparrow \uparrow$ | 0100 | 4 |
| $\vdots$ | $\vdots$ | $\vdots$ |
| $\downarrow \downarrow \downarrow \downarrow$ | 1111 | 15 |

- Convert binary string $\left\{b_{L-1}, \cdots, b_{1}, b_{0}\right\}$ to an integer $I$ :

$$
I=\sum_{i=0}^{L-1} b_{i} 2^{i}
$$

| Quotient | Binary |
| :---: | :---: |
| $7 / / 2=3$ | 1 |
| $3 / / 2=1$ | 1 |
| $1 / / 2=0$ | 1 |
| $0 / / 2=0$ | 0 |

- Convert integer to $/$ to binary string
- at each step divide I/2: remainder forms the binary digit and update $/$ to the quotient

$$
\begin{aligned}
|7\rangle & =|0 \cdots 0111\rangle \\
& =|\uparrow \cdots \uparrow \downarrow \downarrow \downarrow\rangle
\end{aligned}
$$

- Identify the diagonal and off-diagonal components of the Hamiltonian in the computational basis

$$
H_{\mathrm{diag}}=J \sum_{i=0}^{L-1}\left[\sigma_{i}^{z} \sigma_{i+1}^{z}\right]+\sum_{i} h_{i} \sigma_{i}^{z} \quad H_{\text {off-diag }}=J \sum_{i=0}^{L-1}\left[\sigma_{i}^{\chi} \sigma_{i+1}^{x}+\sigma_{i}^{y} \sigma_{i+1}^{y}\right]=J \sum_{i=0}^{L-1}\left[\sigma_{i}^{+} \sigma_{i+1}^{-}+\sigma_{i}^{-} \sigma_{i+1}^{+}\right]
$$

## Diagonal part:

- Take state $|I\rangle=\left\{b_{L-1}^{(I)}, \cdots, b_{1}^{(I)}, b_{0}^{(I)}\right\}$
- Diagonal element $H_{\| I}=\mathcal{E}_{I}$ : the energy of the classical configuration

$$
\mathcal{E}_{I}=J \sum_{i}(-1)^{b_{i}^{(l)}+b_{i}^{(l)}}+\sum_{i} h_{i}(-1)^{b_{i}^{(l)}}
$$

Off-diagonal part:

- spin-flips, connect different spin configurations
- for each state $|\Lambda\rangle$, locate the states $|K\rangle$ for which $\langle K| H \mid \Lambda \neq 0$
- $H_{\text {off-diag }}$ flips nearest-neighbour anti-parallel spins
- for a given state $|\Lambda\rangle$, go to each site $i$

$$
\operatorname{XOR}\left(b_{i}^{(I)}, b_{i+1}^{(I)}\right)= \begin{cases}1 ; & \text { if antiparallel } \\ 0 ; & \text { if parallel }\end{cases}
$$

- if anti-parallel generate a new state

$$
\left\{\cdots b_{i+1}^{(I)} b_{i}^{(I)} \cdots\right\} \Rightarrow\left\{\cdots b_{i}^{(I)} b_{i+1}^{(I)} \cdots\right\}
$$

- generate the decimal representation of new binary string $\equiv K$ and set $H_{I K}=J$

Hamiltonian:

$$
H=J \sum_{i=0}^{L-1}\left[\sigma_{i}^{\times} \sigma_{i+1}^{x}+\sigma_{i}^{y} \sigma_{i+1}^{y}+\sigma_{i}^{z} \sigma_{i+1}^{z}\right]+\sum_{i} h_{i} \sigma_{i}^{z}
$$

Plot of Hamiltonian matrix for $L=6$


Conservation laws/Symmetries

- Total z-magnetisation $S^{z}=\sum_{i} \sigma_{i}^{z}$ is a conserved quantity:

$$
\left[H, S^{z}\right]=0
$$

- Reorder the basis states to group them according to their $S^{z}$
- Hamiltonian becomes block diagonal

Hamiltonian in block diagonal form


- Each block can be diagonalised separately


## Why do we want to do this?

- Often, interested in problems with a fixed magnetisation, total particle number etc.
- If we do not treat each block separately, we can get spurious errors due to accidental near degenracies of eigenvalues, eigenstates mixed between different sectors
- Constructing the full matrix and slicing off the sector of interest is redundant
- Wastage of computational resources; can make the difference between being able and not being able to simulate beyond certain sizes


## Comes at a cost

- We need to work a bit harder to efficiently enumerate the basis with symmetries


## Consider the sector with $S^{z}=0$

Table of configurations and their decimal represetnations

| Configuration | $I$ | $J(I)$ | Configuration | $I$ | $J(I)$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0000111 | 7 | 0 | 100011 | 35 | 10 |
| 001011 | 11 | 1 | 100101 | 37 | 11 |
| 001101 | 13 | 2 | 100110 | 38 | 12 |
| 001110 | 14 | 3 | 101001 | 41 | 13 |
| 010011 | 19 | 4 | 101010 | 42 | 14 |
| 010101 | 21 | 5 | 101100 | 44 | 15 |
| 010110 | 24 | 6 | 110001 | 49 | 16 |
| 011001 | 25 | 7 | 110010 | 50 | 17 |
| 011010 | 26 | 8 | 110100 | 52 | 18 |
| 011100 | 28 | 9 | 111000 | 56 | 19 |

Enumerating the basis

- brute force way is to assign indices in increasing order of the decimal representations of the binary strings
- makes it very inefficient:
- for every decimal $I$, we have to search through the entire (sorted) list to find the appropriate $J(I)$
- two of the possible ways around
- hashing functions
- Lin tables Lin, PRB 1990


## ED of an interacting spin- $1 / 2$ system: construction of basis with symmetries

## Hashing functions

- associate the index of the basis state to the decimal representation of the binray string
- Example:

$$
h(I)=[I(\bmod ) \lambda]+1
$$

Table of configurations and their decimal represetnations

| Configuration | $I$ | $J(I)$ | $h(I)$ | Configuration | $I$ | $J(I)$ | $h(I)$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 000111 | 7 | 0 | 8 | 100011 | 35 | 10 | 13 |
| 001011 | 11 | 1 | 12 | 100101 | 37 | 11 | 15 |
| 001101 | 13 | 2 | 14 | 100110 | 38 | 12 | 16 |
| 001110 | 14 | 3 | 15 | 101001 | 41 | 13 | 19 |
| 010011 | 19 | 4 | 20 | 101010 | 42 | 14 | 20 |
| 010101 | 21 | 5 | 22 | 101100 | 44 | 15 | 22 |
| 010110 | 24 | 6 | 2 | 110001 | 49 | 16 | 4 |
| 011001 | 25 | 7 | 3 | 110010 | 50 | 17 | 5 |
| 011010 | 26 | 8 | 4 | 110100 | 52 | 18 | 7 |
| 011100 | 28 | 9 | 6 | 111000 | 56 | 19 | 11 |

## Hashing functions

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| 000111 | 7 | 0 | 8 | 100011 | 35 | 10 | 13 |
| 001011 | 11 | 1 | 12 | 100101 | 37 | 11 | 15 |
| 001101 | 13 | 2 | 14 | 100110 | 38 | 12 | 16 |
| 001110 | 14 | 3 | 15 | 101001 | 41 | 13 | 19 |
| 010011 | 19 | 4 | 20 | 101010 | 42 | 14 | 20 |
| 010101 | 21 | 5 | 22 | 101100 | 44 | 15 | 22 |
| 010110 | 24 | 6 | 2 | 110001 | 49 | 16 | 4 |
| 011001 | 25 | 7 | 3 | 110010 | 50 | 17 | 5 |
| 011010 | 26 | 8 | 4 | 110100 | 52 | 18 | 7 |
| 011100 | 28 | 9 | 6 | 111000 | 56 | 19 | 11 |

A general issue with hashing functions

- collisions: non-unique hashing values
- very difficult to avoid collisions
- basic idea is to not have a large 1D search
- split the lattice into two parts, $A$ and $B$, and convert the search for the indices into a 2D search
- Define two integers, one for each part

$$
I_{A}=\sum_{i=0}^{L / 2-1} b_{i}^{(I)} 2^{i} \quad I_{B}=\sum_{i=0}^{L / 2-1} b_{i+L / 2}^{(I)} 2^{i}
$$

- Define two vectors $J_{A}\left(I_{A}\right)$ and $J_{B}\left(I_{B}\right)$ so that the position of the configuration represented by integer $I$ is given by

$$
J=J_{A}\left(I_{A}\right)+J_{B}\left(I_{B}\right)
$$

- maximum length of $J_{A}$ and $J_{B}$ is the square root of length of $J(I)$

| Configuration $A$ | $I_{a}$ | $J_{a}\left(I_{a}\right)$ | Configuration $B$ | $I_{b}$ | $J_{b}\left(I_{b}\right)$ | $J=J_{a}+J_{b}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 111 | 7 | 1 | 000 | 0 | 0 | 1 |
| 011 | 3 | 1 | 001 | 1 | 1 | 2 |
| 101 | 5 | 2 | 001 | 1 | 1 | 3 |
| 110 | 6 | 3 | 001 | 1 | 1 | 4 |
| 011 | 3 | 1 | 010 | 2 | 4 | 5 |
| 101 | 5 | 2 | 010 | 2 | 4 | 6 |
| 110 | 6 | 3 | 010 | 2 | 4 | 7 |
| 011 | 3 | 1 | 100 | 4 | 7 | 8 |
| 101 | 5 | 2 | 100 | 4 | 7 | 9 |
| 110 | 6 | 3 | 100 | 4 | 7 | 10 |
| 001 | 1 | 1 | 011 | 3 | 10 | 11 |
| 010 | 2 | 2 | 011 | 3 | 10 | 12 |
| 100 | 4 | 3 | 011 | 3 | 10 | 13 |
| 001 | 1 | 1 | 101 | 5 | 13 | 14 |
| 010 | 2 | 2 | 101 | 5 | 13 | 15 |
| 100 | 4 | 3 | 101 | 5 | 13 | 16 |
| 001 | 1 | 1 | 110 | 6 | 16 | 17 |
| 010 | 2 | 2 | 110 | 6 | 16 | 18 |
| 100 | 4 | 3 | 110 | 6 | 16 | 19 |
| 000 | 0 | 1 | 111 | 7 | 19 | 20 |

